

## CS 170 HW 4

**Due on 2018-09-23, at 9:59 pm**

### 1 (★) Study Group

List the names and SIDs of the members in your study group.

### 2 (★★) Maximum Subarray Sum

Given an array of  $n$  integers, the maximum subarray is the contiguous subarray (potentially empty) with the largest sum. Design a linear algorithm to find the sum of the maximum subarray. For example the maximum subarray of  $[-2, 1, -3, 4, -1, 2, 1, -5, 4]$  is  $[4, -1, 2, 1]$ , whose sum is 6.

Please give a three-part solution of the following format:

- Clearly** describe your algorithm. You can include the pseudocode optionally.
- Write a proof of correctness.
- Write a runtime analysis.

### 3 (★) MST Basics

For each of the following statements, either prove or supply a counterexample. Always assume  $G = (V, E)$  is undirected and connected. Do not assume the edge weights are distinct unless specifically stated.

- Let  $e$  be any edge of minimum weight in  $G$ . Then  $e$  must be part of some MST.
- If  $e$  is part of some MST of  $G$ , then it must be a lightest edge across some cut of  $G$ .
- If  $G$  has a cycle with a unique lightest edge  $e$ , then  $e$  must be part of every MST.
- For any  $r > 0$ , define an  $r$ -path to be a path whose edges all have weight less than  $r$ . If  $G$  contains an  $r$ -path from  $s$  to  $t$ , then every MST of  $G$  must also contain an  $r$ -path from  $s$  to  $t$ .

### 4 (★) Prim's Algorithm

A popular alternative to Kruskal's algorithm is Prim's algorithm, in which the intermediate set of edges  $X$  always forms a subtree, and  $S$  is chosen to be the set of this tree's vertices. We can think of Prim's algorithm as greedily processing one vertex at a time, adding it to  $S$ . The pseudocode below gives the basic outline of Prim's algorithm. See the book for a detailed example of a run of the algorithm.

$S = \{v\}$

$X = \{\}$

While  $S \neq V$ :

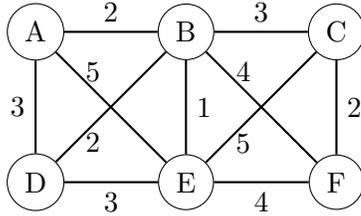
Choose  $t \in V \setminus S, s \in S$  such that  $weight(s,t)$  is minimized

$X = X \cup \{(s,t)\}$

$S = S \cup \{t\}$

Return  $X$

- (a) Run Prim's algorithm on the following graph, starting from A, stating which node you processed and which edge you added at each step .



- (b) Prim's algorithm is very similar to Dijkstra's in that a vertex is processed at each step which minimizes some cost function. These algorithms also produce similar outputs: the union of all shortest paths produced by a run of Dijkstra's algorithm forms a tree. However, the trees they produce aren't optimizing for the same thing. To see this, give an example of a graph for which different trees are produced by running Prim's algorithm and Dijkstra's algorithm. In other words, give a graph where there is a shortest path from a start vertex  $A$  using at least one edge that doesn't appear in any MST.

## 5 (★) Divide and Conquer for MST?

Is the following algorithm correct? If so, prove it. Otherwise, give a counterexample and explain why it doesn't work.

**procedure** FINDMST( $G$ : graph on  $n$  vertices)

If  $n = 1$  return the empty set

$T_1 \leftarrow$  FindMST( $G_1$ : subgraph of  $G$  induced on vertices  $\{1, \dots, n/2\}$ )

$T_2 \leftarrow$  FindMST( $G_2$ : subgraph of  $G$  induced on vertices  $\{n/2 + 1, \dots, n\}$ )

$e \leftarrow$  cheapest edge across the cut  $\{1, \dots, \frac{n}{2}\}$  and  $\{\frac{n}{2} + 1, \dots, n\}$ .

return  $T_1 \cup T_2 \cup \{e\}$ .