

Graphs!

About me

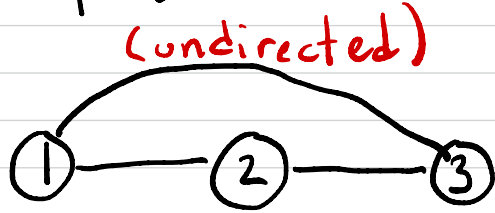
John Wright

4th year at Berkeley

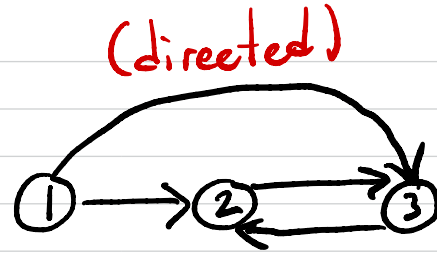
Office: Soda 683

Office hours: Tuesday 3-4pm

Graphs



$$G = (V, E)$$



$$(u, v) \in E \text{ if } u \rightarrow v$$

Parameters: $n = |V|$

$$m = |E| \quad (m \leq n^2)$$



$$\text{deg}(u) = 3$$



$$\begin{aligned} \text{in-deg}(u) &= 1 \\ \text{out-deg}(u) &= 2 \end{aligned}$$



Facebook friend graph
(undirected)

- $n = 2.91$ billion users
- avg user has 338 friends
- $m = 500$ billion edges

Question: Who are my friends?

Best 6 min 15 min 6 min

Yifang Taiwan Fruit Tea, 2516 Bancroft Way

Soda Hall, Berkeley, CA 94709

Add destination

Options

Send directions to your phone Copy link

via Sather Rd 15 min 0.6 mile
Details

via Eshleman Rd 15 min 0.6 mile

via Barrow Ln 14 min 0.6 mile

All routes are mostly flat

Search along the route Restaurants Coffee Groceries Things to do

15 min 0.6 miles

14 min 0.6 miles

15 min 0.6 miles

University of California, Berkeley

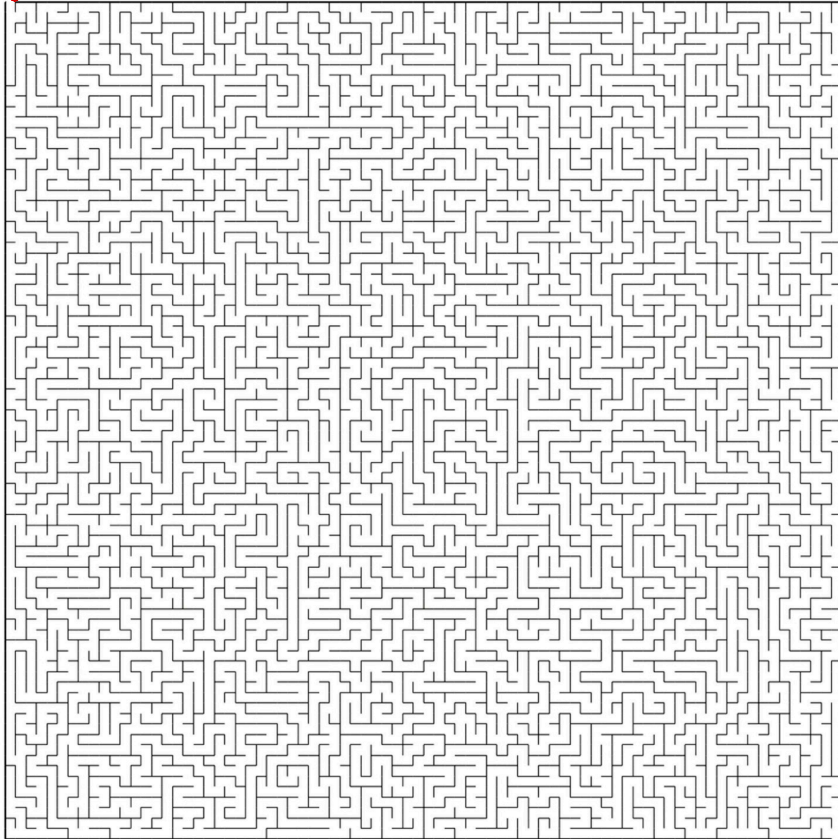
Yifang Taiwan Fruit Tea

Soda Hall

Google Maps

Q: Shortest path to boba?

entrance u



exit v

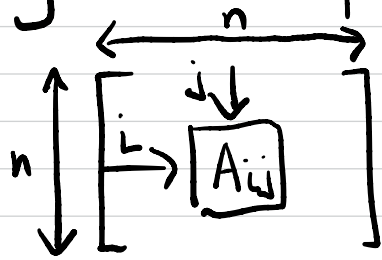
Maze solving

Q: Is u connected
to v ?

Representing graphs on computers

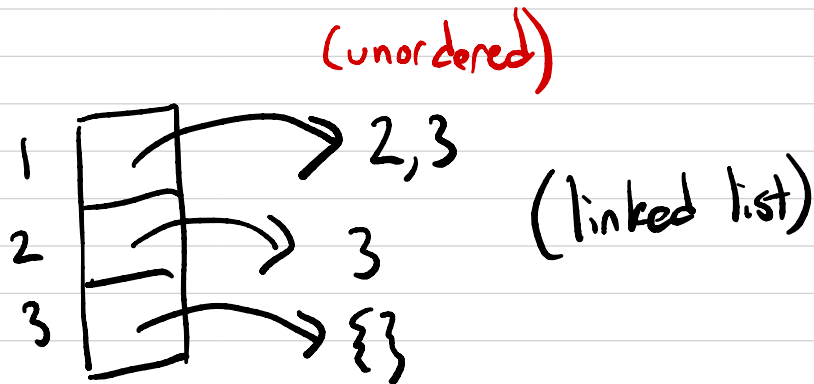
$$V = \{1, \dots, n\}$$

(1) adjacency matrix representation



$$A_{ij} = \begin{cases} 1 & \text{if } (i, j) \in E \\ 0 & \text{o.w.} \end{cases}$$

(2) adjacency list representation



	Adjacency matrix	Adjacency list
size		
time to: { answer		
{ is $(u,v) \in E$?		
{ enumerate		
{ u 's neighbors		

Facebook graph w/ a adjacency matrix

$$\text{size } n^2 = (2.91 \text{ billion})^2 \text{ space}$$

$$= 8.5 \times 10^{18} = 1 \text{ million TB}$$

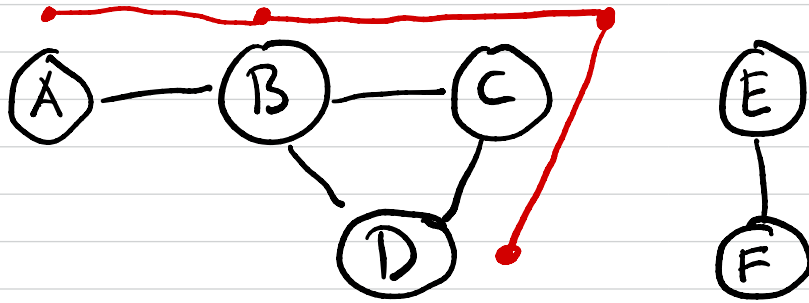
$$\text{computing friends list} = 2.91 \text{ billion steps}$$

Graph Exploration

Useful for:

1. Is there a path from u to v ?
2. Is G connected?
3. What are G 's connected components?

Connectivity for undirected graphs (\exists path from u to v ?)



explore(G, u)
visited[u] = true

for v s.t. $(u, v) \in E$
if visited[v] = false
explore(G, v)

boolean array visited[n]
(init to all 0's)

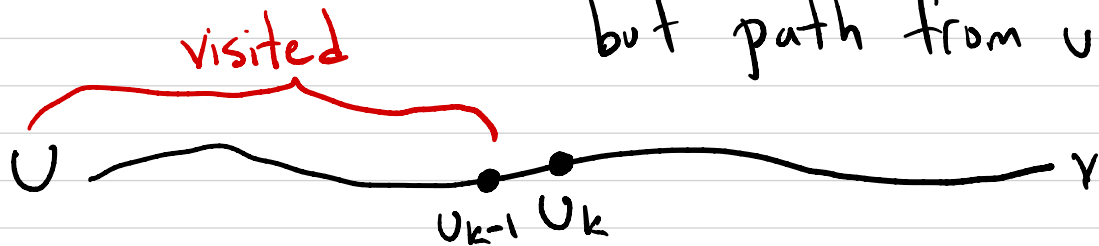
Property: $\text{explore}(G, u)$ visits exactly the vertices v
s.t. G has path from u to v

Pf: 1. v explored \Rightarrow path from u to v

2. path from u to $v \Rightarrow$ explore v

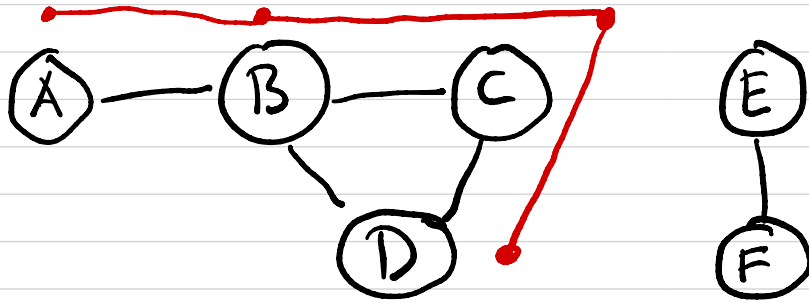
Assume false, so v not explored

but path from u to v



Should call $\text{explore}(G, u_k)$. Contradiction!

Connectivity for undirected graphs (\exists path from u to v ?)



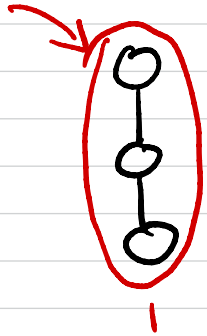
explore(G, u)
visited[u] = true

for v s.t. $(u, v) \in E$
if visited[v] = false
explore(G, v)

dfs(G)
boolean array visited [n]
(init to all 0's)

for $v \in V$
if visited[v] = false
explore(G, v)

Computing G 's connected components (undirected)



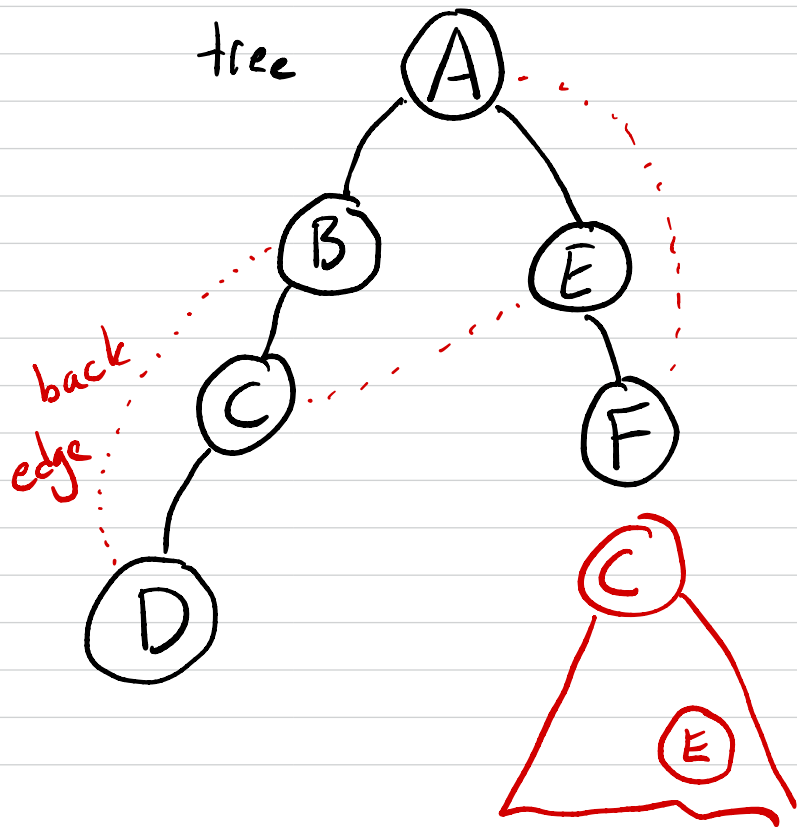
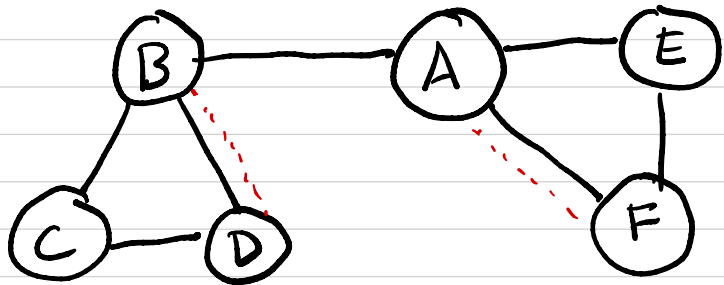
(3 connected components)

```
explore(G, u)
  visited[u] = true
  ccnum[u] = count
```

```
for v s.t. (u, v) ∈ E
  if visited[v] = false
    explore(G, v)
```

```
dfs(G)
  boolean array visited[n]
  (init to all 0's)
  count = 1
  int array ccnum[n]
  for v ∈ V
    if visited[v] = false
      explore(G, v)
      count = count + 1
```

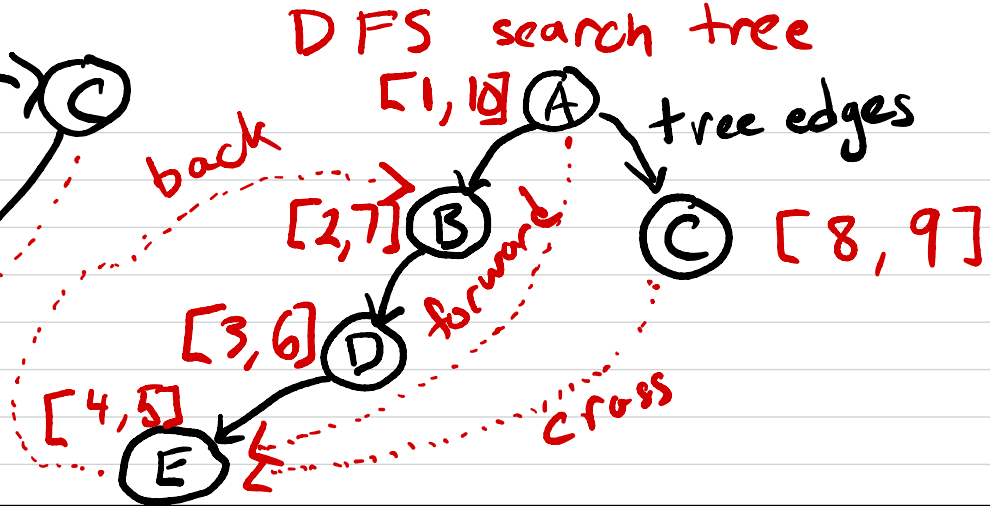
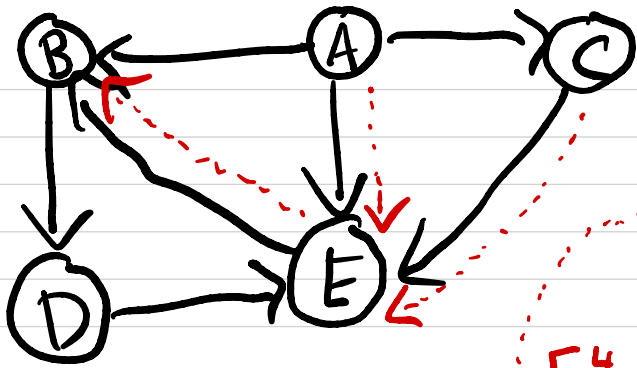
The DFS tree



Runtime of DFS

Only call $\text{explore}(G, v)$ exactly once, for each v

Runtime of $\text{explore}(G, v)$:



```

explore (G, u)
  visited[u] = true
  pre[u] = clock
  clock = clock + 1
  for v s.t. (u, v) ∈ E
    if visited[v] = false
      explore(G, v)
  post[u] = clock
  clock = clock + 1
  
```

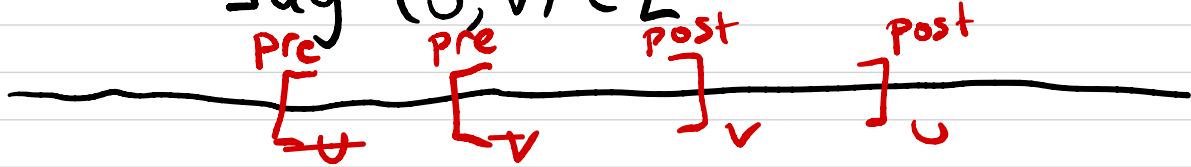
```

dfs (G)
  boolean array visited [n]
  (init to all 0's)
  clock = 1
  int array pre [n], post [n]
  for v ∈ V
    if visited[v] = false
      explore(G, v)
  
```

Classifying edges using pre & post #'s

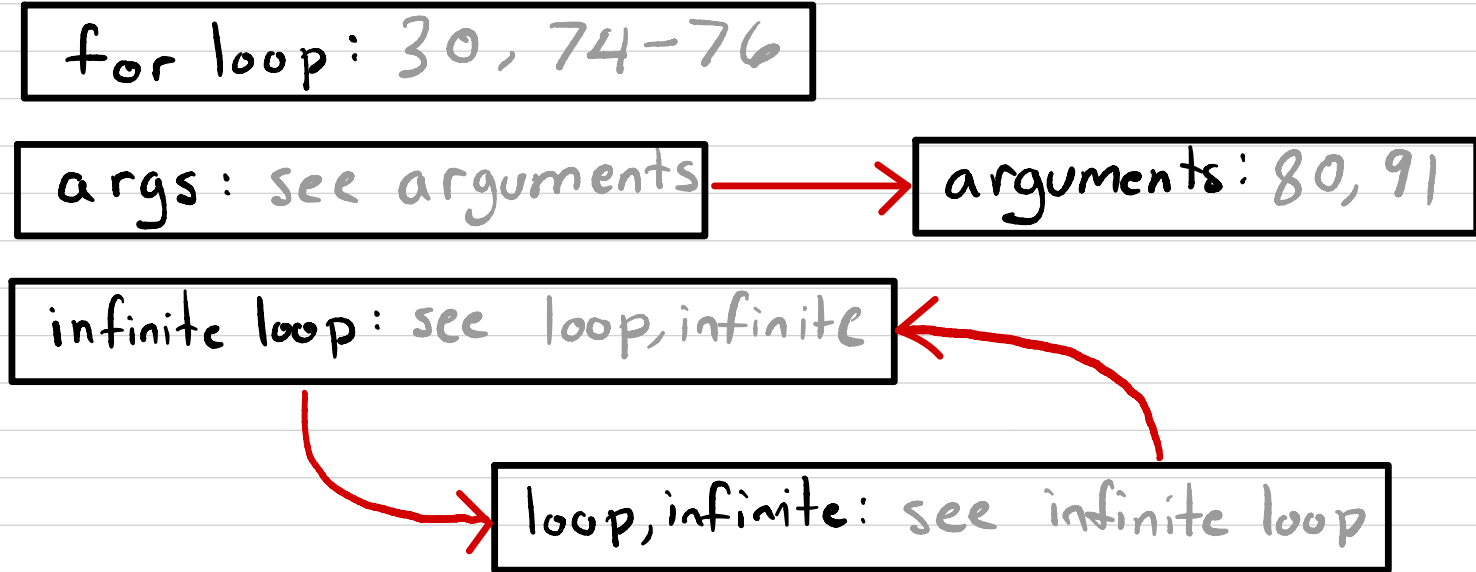
Say $(u, v) \in E$

Clock



Application # 1: Cycle detection

Book index:

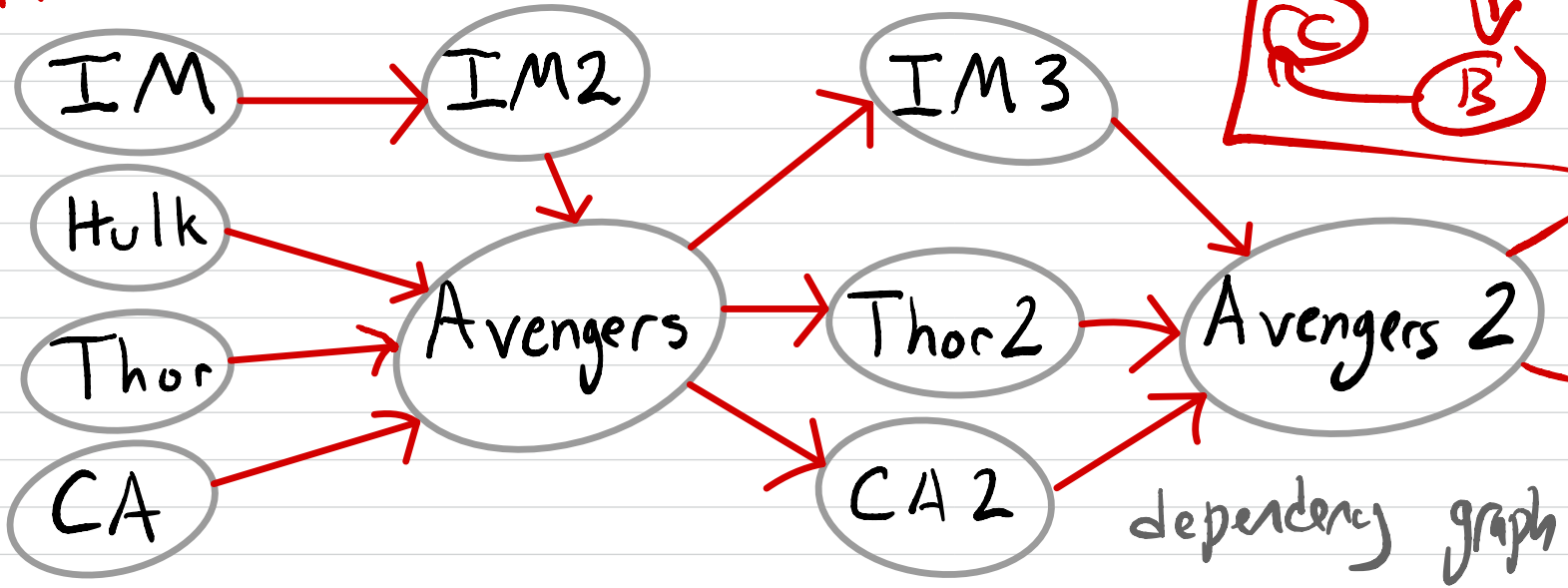


Q: Does my graph have a cycle?

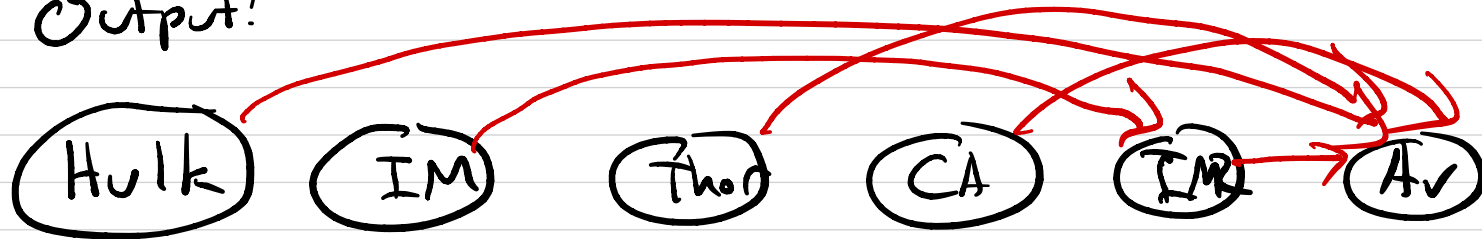
Def: A directed acyclic graph (DAG)
is a directed graph w/ no cycles.

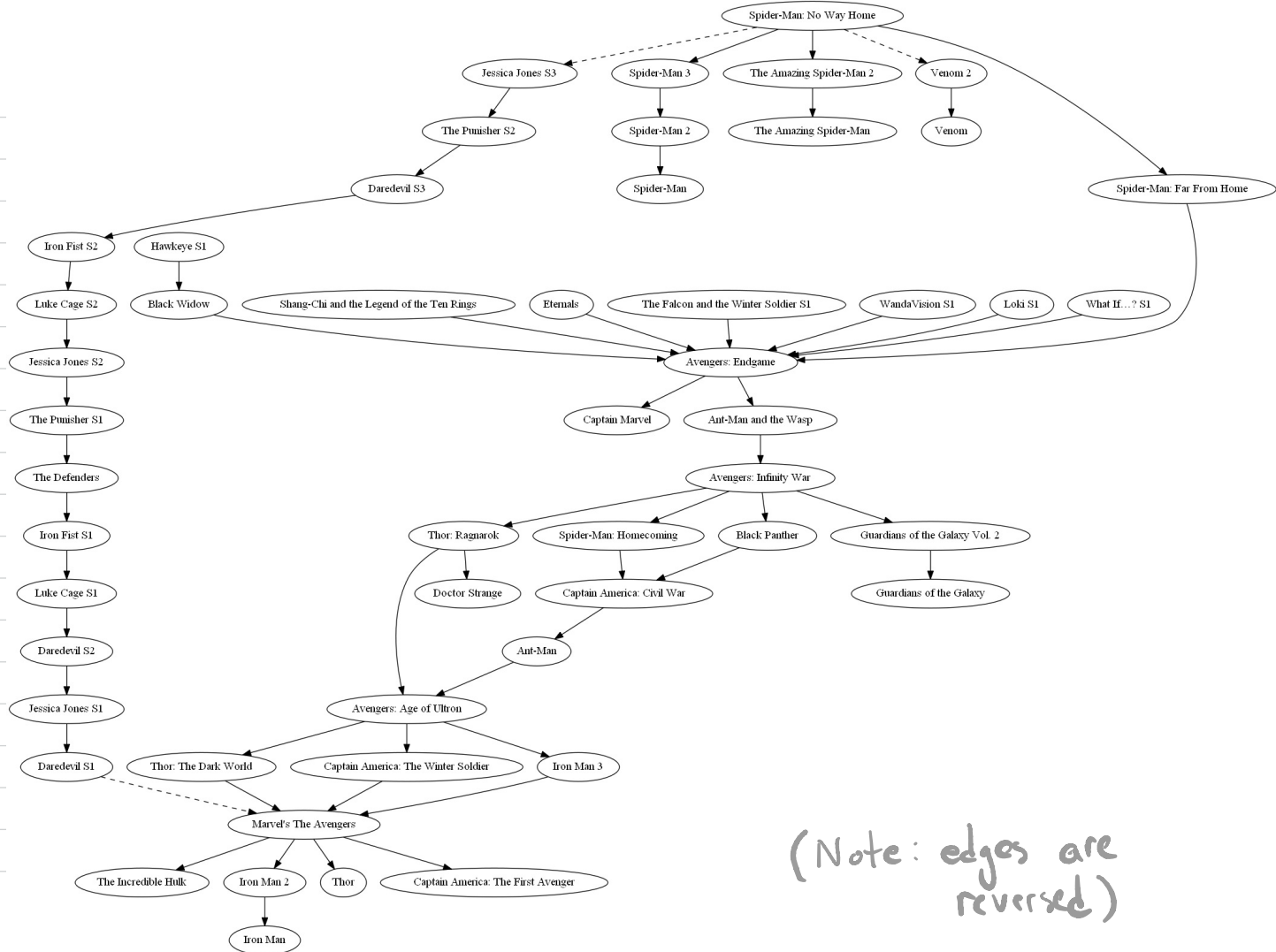
Claim: Suppose we run DFS on G .
Then G is a DAG iff

Application #2: Topological Sort



Output:





(Note: edges are reversed)

Topological sort

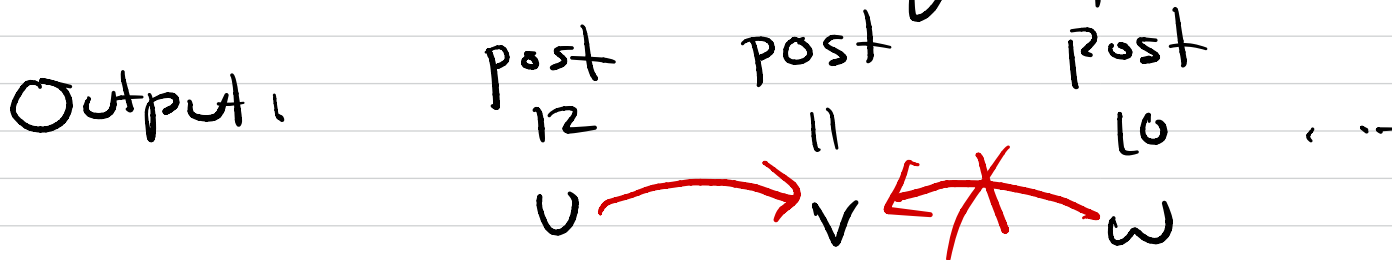
Input: DAG $G = (V, E)$

Output: Ordering of vertices v_1, \dots, v_n
s.t. all edges go left \rightarrow right

Claim: Suppose we DFS on G .
Then for all $(u, v) \in E$, $\text{post}(u) > \text{post}(v)$.

Pf: G is a DAG \Rightarrow no back edges
 $\Rightarrow \text{post}(u) > \text{post}(v)$ for all edges. \square

Alg: Run DFS on G .
Sort vertices from highest post to lowest.



Application #3: Connected components in digraphs

