1 Streaming for Voting

Consider the following scenario. Votes are being cast for a major election, but due to a lack of resources, only one computer is available to count the votes. Furthermore, this computer only has enough space to store one vote at a time, plus a single extra integer. Each vote is a single integer 0 or 1, denoting a vote for Candidate A and Candidate B respectively.

(a) Come up with an algorithm to determine whether candidate A or B won, or if there was a tie.

(b) Consider now an election with 3 candidates. Say there is a winner only if a candidate recieves more than 50 percent of the vote, otherwise there is no winner. If we’re given another integer’s worth of storage, come up with an algorithm to determine the winner if there is one. For simplicity, your algorithm can output any of the candidates in the case that there is no winner (not necessarily the one with the most votes). Votes are now numbered 0, 1, 2.

2 Universal Hashing

Let $[m]$ denote the set $\{0, 1, ..., m - 1\}$. For each of the following families of hash functions, determine whether or not it is pairwise-independent. If it is universal, determine how many random bits are needed to choose a function from the family.

(a) $H = \{ h_{a_1, a_2} : a_1, a_2 \in [m] \}$, where $m$ is a fixed prime and

$$h_{a_1, a_2}(x_1, x_2) = a_1 x_1 + a_2 x_2 \mod m$$

Notice that each of these functions has signature $h_{a_1, a_2} : [m]^2 \to [m]$ , that is, it maps a pair of integers in $[m]$ to a single integer in $[m]$.

(b) $H$ is as before, except that now $m = 2^k$ is some fixed power of 2.

(c) $H$ is the set of all functions $f : [m] \to [m - 1]$. 
